

These notes are from <http://math.cmu.edu/~coco/teaching/discrete20/rec5.pdf>

This is a review day, so hopefully you came prepared with your own questions! Here are a couple additional problems to think about (solutions are on the next page).

Problem 1. Prove that $\sum_{i=0}^k \binom{m}{i} \binom{n}{k-i} = \binom{m+n}{k}$ in two different ways:

1. Using double counting.
2. Using the binomial theorem.

Problem 2. Fix $k \geq 2$ and let $a_k(n)$ denote the number of words in $[k]^n$ which have an even number of 1's. Find a formula for $a_k(n)$.

(Hint: It may be helpful to consider also $b_k(n)$ to be the number of words in $[k]^n$ which have an odd number of 1's.)

Problem 3. Let Ω be a finite set and let $g: 2^\Omega \rightarrow \mathbb{R}$ be any function. Define the function $f: 2^\Omega \rightarrow \mathbb{R}$ by

$$f(S) = \sum_{T \subseteq S} g(T),$$

for all $S \subseteq \Omega$. Prove that for any $S \subseteq \Omega$,

$$g(S) = \sum_{T \subseteq S} (-1)^{|S|-|T|} f(T).$$

(This is known as the Möbius inversion formula on the Boolean lattice.)

Solution to Problem 1. Part 1 was actually Problem 6 on Homework 1, so please review that solution.

For Part 2, we'll use the binomial theorem. Firstly,

$$(1+x)^{m+n} = \sum_{k=0}^{m+n} \binom{m+n}{k} x^k.$$

On the other hand,

$$\begin{aligned} (1+x)^{m+n} &= (1+x)^m(1+x)^n = \left(\sum_{k=0}^m \binom{m}{k} x^k \right) \left(\sum_{k=0}^n \binom{n}{k} x^k \right) \\ &= \sum_{k=0}^{m+n} \left(\sum_{i=0}^k \binom{m}{i} \binom{n}{k-i} \right) x^k, \end{aligned}$$

so

$$\sum_{k=0}^{m+n} \binom{m+n}{k} x^k = \sum_{k=0}^{m+n} \left(\sum_{i=0}^k \binom{m}{i} \binom{n}{k-i} \right) x^k.$$

Since this holds for every x , it must be the case that for every k , the coefficient of x^k is the same on both sides; therefore, $\binom{m+n}{k} = \sum_{i=0}^k \binom{m}{i} \binom{n}{k-i}$. \square

Solution to Problem 2. Let $A_k(n)$ denote the set of words in $[k]^n$ with an even number of 1's and let $B_k(n)$ denote the set of words in $[k]^n$ with an odd number of 1's; therefore $a_k(n) = |A_k(n)|$ and $b_k(n) = |B_k(n)|$. Observe that $A_k(n), B_k(n)$ forms a partition of $[k]^n$ so $a_k(n) + b_k(n) = k^n$. We will show that $a_k(n) - b_k(n) = (k-2)^n$, which will imply that

$$a_k(n) = \frac{1}{2}(k^n + (k-2)^n).$$

For a word $w \in [k]^n$, let $z(w)$ denote the number of 1's in w . Define the sign of w to be $\sigma(w) = (-1)^{z(w)}$ and set $Q := \sum_{w \in [k]^n} \sigma(w)$. We first observe that

$$Q = \sum_{w \in A_k(n)} (-1)^{\text{even}} + \sum_{w \in B_k(n)} (-1)^{\text{odd}} = a_k(n) - b_k(n).$$

Now, for $w \in [k]^n$, define $F(w) = \min\{i \in [n] : w_i \in \{1, 2\}\}$; that is the smallest coordinate of w which is either a 1 or a 2. Observe that $F(w)$ is undefined if and only if w has neither a 1 nor a 2; thus $\{w \in [k]^n : F(w) \text{ undefined}\} = \{3, \dots, k\}^n$, which has size $(k-2)^n$. Furthermore, if $F(w)$ is undefined then $\sigma(w) = 1$ since w has no 1's.

Now, consider w for which $F(w)$ is defined and let $f(w)$ denote the word in which the $F(w)$ 'th coordinate is flipped from a 1 to a 2 or a 2 to a 1. For instance, $f(342122) = 341122$. Observe that f is an involution on those w 's for which $F(w)$ is defined; furthermore the number of 1's in w and $f(w)$ differ by exactly one, so $\sigma(w) = -\sigma(f(w))$.

We can therefore compute

$$Q = \sum_{\substack{w: \\ F(w) \text{ undefined}}} \sigma(w) + \sum_{\{w, f(w)\}} (\sigma(w) + \sigma(f(w))) = (k-2)^n. \quad \square$$

Solution to Problem 3. We first show that the claimed g is valid.

$$\begin{aligned}
\sum_{T \subseteq S} g(T) &= \sum_{T \subseteq S} \sum_{R \subseteq T} (-1)^{|T|-|R|} f(R) \\
&= \sum_{R \subseteq S} \sum_{T: R \subseteq T \subseteq S} (-1)^{|T|-|R|} f(R) \\
&= \sum_{R \subseteq S} f(R) \left(\sum_{T: R \subseteq T \subseteq S} (-1)^{|T|-|R|} \right) \\
&= \sum_{R \subseteq S} f(R) \cdot \mathbf{1}[R = S] \\
&= f(S).
\end{aligned}$$

Of course, just because the claimed g satisfies the formula, doesn't mean that we're done with the problem. Recall that g was some fixed function in the problem statement and f was defined from g . So far, we've shown only that the claimed formula will define the same f , not that it was actually the function we started with. In other words, we need to show that this is the *only* g that will work.

To this end, suppose that $g, h: 2^\Omega \rightarrow \mathbb{R}$ are such that for every $S \subseteq \Omega$,

$$\sum_{T \subseteq S} g(T) = \sum_{T \subseteq S} h(T);$$

we need to show that $g = h$. Suppose not, then $g(S) \neq h(S)$ for some $S \subseteq \Omega$. Since Ω is a finite set, we may consider the smallest S for which $g(S) \neq h(S)$ (here, smallest simply means smallest size). Note that there may be multiple S 's of smallest size; however, we know that $g(T) = h(T)$ for all $|T| < |S|$, which is all that is important. We then compute

$$\begin{aligned}
\sum_{T \subseteq S} g(T) &= g(S) + \sum_{T \subsetneq S} g(T) = g(S) + \sum_{T \subsetneq S} h(T) \\
&\neq h(S) + \sum_{T \subsetneq S} h(T) = \sum_{T \subseteq S} h(T);
\end{aligned}$$

contradicting our original assumption. □