

These notes are from <http://math.cmu.edu/~cocox/teaching/discrete20/rec12.pdf>

Let  $G$  be a graph and let  $\chi: E \rightarrow \{0, 1\}$  be a 2-coloring of the edges of  $G$ . In this coloring, we can define  $\deg_i^\chi(v) = |\{u \in V : \chi(uv) = i\}|$  for  $i \in \{0, 1\}$ ; that is,  $\deg_i^\chi(v)$  is the number of edges of color  $i$  incident to  $v$  in the coloring  $\chi$ . For  $\chi: E \rightarrow \{0, 1\}$ , we define the *discrepancy* of  $\chi$  to be

$$\text{disc}(G, \chi) \stackrel{\text{def}}{=} \max_{v \in V} |\deg_0^\chi(v) - \deg_1^\chi(v)|.$$

Of course, for any graph  $G$  and 2-coloring  $\chi$ , we have  $\text{disc}(G, \chi) \leq \Delta(G)$ , which is tight when  $\chi$  gives every edge the same color. So the interesting question is to try to minimize the discrepancy; hence, define

$$\text{disc}(G) \stackrel{\text{def}}{=} \min_{\chi: E \rightarrow \{0, 1\}} \text{disc}(G, \chi).$$

Observe that if  $G$  has a vertex of odd degree, then  $\text{disc}(G) \geq 1$ . Moreover, observe that if  $G$  is a cycle with an odd number of vertices, then no matter how the edges are 2-colored, there must be some vertex incident to edges of only a single color, so  $\text{disc}(G) = 2$ . It turns out that an odd cycle is a worst-case scenario (see the exercise below for a classification of all worst-case scenarios).

**Claim 1.** *For any graph  $G$ ,  $\text{disc}(G) \leq 2$ .*

*Proof.* We will complete the proof in a couple steps.

**Step 1:  $G$  is connected and  $\deg(v)$  is even for every  $v \in V$ .** Here we prove something slightly stronger: For any fixed  $v \in V$ , there is a coloring  $\chi: E \rightarrow \{0, 1\}$  such that  $|\deg_0^\chi(v) - \deg_1^\chi(v)| \leq 2$  and  $\deg_0^\chi(u) = \deg_1^\chi(u)$  for all  $u \neq v$ .

Since  $G$  is connected and every vertex has even degree, we know that  $G$  has an Eulerian walk. We may suppose that this walk starts at  $v$ : label the edges  $e_1, e_2, \dots, e_m$  in the order they're traversed and define  $\chi(e_i) = i \bmod 2$ . Consider any vertex  $u \neq v$ . Observe that if  $u \in e_i$ , then either  $u \in e_{i-1}$  or  $u \in e_{i+1}$  (but not both). Hence, we can pair up edges of opposite color incident to  $u$ , implying that  $\deg_0^\chi(u) = \deg_1^\chi(u)$ . The same logic almost works for  $v$ , except for the fact that  $e_1, e_m$  are paired up and these edges may have the same color. Hence, we have  $\deg_0^\chi(v) = \deg_1^\chi(v)$  if  $m$  is even and  $\deg_1^\chi(v) = \deg_0^\chi(v) + 2$  if  $m$  is odd.<sup>1</sup>

**Step 2:  $G$  is connected and some vertex has odd degree.** Let  $U \subseteq V$  denote the set of vertices of  $G$  with odd degree. We build a new graph  $G'$  from  $G$  by adding a new vertex which is connected to each vertex of  $U$ . Formally,  $V(G') = V \cup \{v'\}$  and  $E(G') = E \cup \{\{v', u\} : u \in U\}$ . The handshaking lemma tells us that  $|U|$  is even, so since  $|U| \geq 1$  by assumption, we know that  $G'$  is connected and every vertex of  $G'$  has even degree.

By step 1, we may find a coloring  $\chi': E(G') \rightarrow \{0, 1\}$  such that  $|\deg_0^{\chi'}(v') - \deg_1^{\chi'}(v')| \leq 2$  and  $\deg_0^{\chi'}(v) = \deg_1^{\chi'}(v)$  for all  $v \in V(G') \setminus \{v'\} = V$ . Thus, by deleting  $v'$  and its incident edges, we are left with a coloring  $\chi: E \rightarrow \{0, 1\}$  for which  $|\deg_0^\chi(v) - \deg_1^\chi(v)| \leq 1$  for all  $v \in V$ .<sup>2</sup>

<sup>1</sup>Observe that we've proved something more than claimed here: If  $G$  is a connected graph with an even number of edges and every vertex has even degree, then  $\text{disc}(G) = 0$ .

<sup>2</sup>Observe that we've proved something more than claimed here: If  $G$  is a connected graph with at least one vertex of odd degree, then  $\text{disc}(G) = 1$ .

**Step 3:  $G$  is arbitrary.** We can decompose  $G$  into  $G = G_1 \cup G_2 \cup \cdots \cup G_k$  where the  $G_i$ 's are vertex disjoint and each  $G_i$  is connected. Observe that since the  $G_i$ 's are vertex disjoint,  $\text{disc}(G) = \max_{i \in [k]} \text{disc}(G_i)$ . Putting together steps 1 and 2, since each  $G_i$  is connected,  $\text{disc}(G_i) \leq 2$ , and so  $\text{disc}(G) \leq 2$  as well.  $\square$

**Exercise.** Let  $G$  be a connected graph. Show that  $\text{disc}(G) = 2$  if and only if  $G$  has an odd number of edges and every vertex has even degree.

**Exercise.** Let  $G$  be a connected graph and suppose that  $\chi: E \rightarrow \{0, 1\}$  is any 2-coloring for which  $\text{disc}(G, \chi) = 0$ . Show that  $G$  has an Eulerian walk  $e_1, \dots, e_m$  such that  $\chi(e_i) = i \pmod 2$ .